Dr. Marques Sophie Office 519 Number theory

Fall Semester 2013 marques@cims.nyu.edu

# Problem Set #8

## Due monday november 4th in Class

**Exercise 1:**  $(\star)$  4 points

Let  $\phi$  be Euler's function. Show that  $\phi(n^m) = n^{m-1}\phi(n)$  for all natural number n, m. Solution:

Let  $n = p_1^{a_1} \dots p_r^{a_r}$  be the prime decomposition for n. Since  $\phi$  is multiplicative on coprime elements, we have that  $\phi(n^m) = \phi(p_1^{ma^1}) \dots \phi(p_r^{ma_r})$ . Next, we know that if q is prime and  $k \ge 1$  then  $\phi(q^k) = q^{k-1}(q-1)$ . Written another way reads  $\phi(q^k) = q^{k-1}(q-1)$ . Using this with what we have already written we get

$$\phi(n^m) = p_1^{ma_1 - 1} \dots p_r^{ma_r - 1} \phi(p_1) \dots \phi(p_r) = \frac{p_1^{ma_1} \dots p_r^{ma_r}}{p_1 \dots p_r} \phi(n) = n^{m-1} \phi(n)$$

where we have used the multiplicativity of  $\phi$  on the primes to get the second equality.

#### **Exercise 2:** $(\star)$ 4 points

Let N be a product of two distinct primes. Show that if we know  $\varphi(N)$ , then we can easily factorize N.

#### Solution:

Let N = pq where p < q are distinct primes. Then,

$$\varphi(N) = (p-1)(q-1) = pq - (p+q) + 1$$

Define  $K := N - \varphi(N) + 1 = p + q$ . Note that we can compute K from N and  $\varphi(N)$  without knowing p and q. We have

$$(x-p)(x-q) = x^{2} - (p+q)x + pq = x^{2} - Kx + N$$

Thus, by solving the roots of this quadratic polynomial, we obtain that

$$p = \frac{K - \sqrt{K^2 - 4N}}{2} \qquad \qquad q = \frac{K + \sqrt{K^2 - 4N}}{2}$$

In this way, we can obtain prime factorization of N quickly.

#### **Exercise 3:** $(\star)$ 4 points

Find the last two digits of the decimal expansion of  $3^{1123}$  (For example the last two digits of 1729 are 29).

## Solution:

Notice that finding the last two digits of the  $3^{1123}$  is equal to finding w in the equation:  $3^{1123} \equiv x \mod 100$ , notice that gcd(3, 100) = 1 so, as it follows from Euler's theorem  $3^{\phi(100)} \equiv 1 \mod 100$ , we will find  $\phi(100)$ :  $100 = 2^2 \times 5^2$ , therefore  $\phi(100) = (2-1) \times 2^{2-1} \times (5-1) \times 5^{2-1} = 40$ , notice that  $1112 = 28 \times 40 + 3$ , therefore  $3^{1123} = (3^{40})^{28} \times 3^3 \equiv 1 \times 3^3 \mod 100 = 27 \mod 100$ , then x = 27, so the last two digits of  $3^{1123}$  are 27.

**Exercise 4:**  $(\star)$  4 points

Show that there is no solution to the equation  $\phi(n) = 14$ .

## Solution:

Assume that  $n = p_1^{t_1} p_2^{t_2} \dots p_k^{t_k}$  is the prime factorization of n. Using the the formula derived in class, if then  $\phi(n) = 14$  means that

$$14 = p_1^{t_1-1}(p_1-1)p_2^{t_2-1}(p_2-1)\dots p_k^{t_k-1}(p_k-1)$$

On account of the  $p_i - 1$  terms in the expression, the only primes  $p_i$  in n must be such that  $p_i - 1$  divides 14. Thus the only primes possible in n are 2 and 3. This means that  $n = 2^a 3^b$  where  $0 \le a, b$  so that

$$14 = 2^{a-1} \times 1 \times 3^{b-1} \times 2$$

where the  $p^{j-1}(p-1)$  term is present only if the p occurs in the prime factorization of n.

Note that if b > 1 then 3|14 a contradiction thus b = 0, 1. But if b = 1 then the equation becomes  $14 = 2^{a-1}3^{1-1}2$  which implies that  $2^{a-1} = 7$  which is impossible. Thus b = 0. But if b = 0 then then the equation becomes  $14 = 2^{a-1}$  which once again is impossible. Thus there can be no solution to the equation  $\phi(n) = 12$ .

## **Exercise 5:** $(\star)$ 4 points

You receive a message that was encrypted using the RSA system with public key (65, 29), where 65 is the base and 29 is the exponent. The encrypted message in two blocks, is 3/2. Find the private key and decrypt the message.

## Solution:

First we find  $\phi(65)$ . The prime factorization of 65 is  $5 \times 13$ , hence  $\phi(65) = \phi(5)\phi(13) = (5-1)(13-1) = 48$ . To find  $\beta$ , we can apply the Euclidean algorithm to 29 and 48 and we find  $5 \times 29 - 3 \times 48 = 1$  which implies that 5 is the inverse of 29 modulo 48. Now that we know the private key, the decrypted message is  $b_1/b_2$ , where  $b_1 \equiv 3^5 \mod 65$ ,  $b_2 \equiv 2^5 \mod 65$ , and  $0 \le b_1, b_2 < 65$ . We find that  $b_1 = 48$ ,  $b_2 = 32$ .

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 $<sup>^{1}(\</sup>star) = easy , (\star\star) = medium, (\star\star\star) = challenge$